

## Lecture 2

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# 1 Duality Theory

## 1.1 Motivation Examples

**Example 1.1.** Let us consider the following optimization problem:

$$\begin{aligned} \min_{\mathbf{x}} f(\mathbf{x}), \\ \text{s.t. } c_i(\mathbf{x}) = 0, i = 1, \dots, m. \end{aligned}$$

If  $c_i(\mathbf{x}) = \mathbf{a}_i^\top \mathbf{x} - b_i$ , we have the optimality condition for constrains  $A\mathbf{x} = \mathbf{b}$ . If  $c_i$  is not a linear function, the optimality condition is  $\langle \nabla f(\mathbf{x}^*), \mathbf{y} - \mathbf{x}^* \rangle \geq 0$ , for all  $\mathbf{y} \in \mathcal{X} = \{\mathbf{x} | c_i(\mathbf{x}) = 0, i = 1, \dots, m\}$ . This means we have no an equation system to solve the optimal point compared with the equality constrains.

**Example 1.2.** LAD regression:  $\min_{\mathbf{x}} \|A\mathbf{x} - \mathbf{b}\|_1$ .

- Sub-gradient descent:  $\mathbf{x}^{t+1} = \mathbf{x}^t - s_t \partial_{\|\cdot\|_1}(A\mathbf{x}^t - \mathbf{b})$ . The speed is  $O(\frac{1}{\sqrt{T}})$ .
- Proximal Gradient Descent: consider  $\min_{\mathbf{x}} f(\mathbf{x}) + \|A\mathbf{x} - \mathbf{b}\|_1$ , where  $f(\mathbf{x}) = 0$ . Then the corresponding PGD algorithm is

$$\begin{cases} \mathbf{x}^{t+1} = \text{prox}_{\alpha\|A\mathbf{x}-\mathbf{b}\|_1}(\mathbf{x}^t), \\ \text{prox}_{\alpha\|A\mathbf{x}-\mathbf{b}\|_1}(\mathbf{x}^t) = \arg \min \left\{ \frac{1}{2\alpha} \|\mathbf{x} - \mathbf{x}^t\|^2 + \|A\mathbf{x} - \mathbf{b}\|_1 \right\}. \end{cases}$$

**Example 1.3.** Fused LASSO [?]:

$$\min_{\mathbf{x}} \frac{1}{2} \|A\mathbf{x} - \mathbf{b}\|^2 + \lambda \|F\mathbf{x}\|_1, \tag{1}$$

where  $F \in \mathbb{R}^{(n-1) \times n}$  and

$$F_{ij} = \begin{cases} 1, & j = i + 1, \\ -1, & j = i, \\ 0, & \text{otherwise.} \end{cases}$$

## 1.2 The Lagrange Dual Function

We consider that

$$\begin{aligned} \min_{\mathbf{x}} f_0(\mathbf{x}), \\ \text{s.t. } f_i(\mathbf{x}) \leq 0, i = 1, \dots, m, \\ h_j(\mathbf{x}) = 0, j = 1, \dots, l. \end{aligned}$$

**Definition 1.4.** We define that *Lagrangian*  $L : \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^l \rightarrow \mathbb{R}$  is

$$L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) := f_0(\mathbf{x}) + \sum_{i=1}^m \lambda_i f_i(\mathbf{x}) + \sum_{j=1}^l \nu_j h_j(\mathbf{x}), \quad (2)$$

where  $\boldsymbol{\lambda} = (\lambda_1, \dots, \lambda_m)^\top$  and  $\boldsymbol{\nu} = (\nu_1, \dots, \nu_l)^\top$  are denoted as *dual variables* or *Lagrange multipliers*.

**Definition 1.5.** Define the *Lagrange dual function* as

$$g(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\mathbf{x} \in D} L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}), \quad (3)$$

where  $D = \{\cap_{i=0}^m \text{dom}(f_i)\} \cap \{\cap_{j=1}^l \text{dom}(h_j)\}$ .

**Theorem 1.6.** Let us define that  $p^* = \min_{\mathbf{x} \in \mathcal{X}} f_0(\mathbf{x})$ , then

$$g(\boldsymbol{\lambda}, \boldsymbol{\nu}) \leq p^*$$

for any  $\boldsymbol{\lambda} \geq 0$ .

*Proof.* Suppose that  $\bar{\mathbf{x}} \in \mathcal{X}$ , then  $\sum_{i=1}^m \lambda_i f_i(\bar{\mathbf{x}}) + \sum_{j=1}^l \nu_j h_j(\bar{\mathbf{x}}) \leq 0$ . Thus,

$$\begin{aligned} g(\boldsymbol{\lambda}, \boldsymbol{\nu}) &= \inf_{\mathbf{x} \in D} L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \leq L(\bar{\mathbf{x}}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \\ &= f_0(\bar{\mathbf{x}}) + \sum_{i=1}^m \lambda_i f_i(\bar{\mathbf{x}}) + \sum_{j=1}^l \nu_j h_j(\bar{\mathbf{x}}) \\ &\leq f_0(\bar{\mathbf{x}}), \end{aligned}$$

for all  $\bar{\mathbf{x}} \in \mathcal{X}$ . Therefore,  $g(\boldsymbol{\lambda}, \boldsymbol{\nu}) \leq f_0(\mathbf{x}^*) = p^*$ . ■

**Remark 1.7.** • *Theorem 1.6* shows the Lagrange dual function gives a nontrivial lower bound on  $p^*$  only when  $\boldsymbol{\lambda} \geq 0$  and  $(\boldsymbol{\lambda}, \boldsymbol{\nu}) \in \text{dom}(g)$ . We refer to a pair  $(\boldsymbol{\lambda}, \boldsymbol{\nu}) \in \text{dom}(g)$  with  $\boldsymbol{\lambda} \geq 0$  as *dual feasible variables*.

- $g(\boldsymbol{\lambda}, \boldsymbol{\nu})$  is always concave.

**Definition 1.8.** For each pair  $(\boldsymbol{\lambda}, \boldsymbol{\nu}) \in \text{dom}(g)$  with  $\boldsymbol{\lambda} \geq 0$ , the Lagrange dual function gives us a lower bound of  $p^*$ . A natural question is what is the best lower bound that can be obtained from the Lagrange

dual function. This leads to the following optimization problem:

$$q^* = \max_{\lambda, \nu} g(\lambda, \nu), \quad (4)$$

$$s.t. \lambda \succeq 0. \quad (5)$$

The previous problem is called *Lagrange dual problem* and  $(\lambda^*, \nu^*)$  are the *dual optimal variables* or *optimal Lagrange multipliers*.

The Lagrange dual problem is a convex optimization since the objective to be maximized is concave and the constraint is convex, whether or not the primal problem is convex.

**Definition 1.9. Weak Duality:**  $q^* \leq p^*$ .

**Strong Duality:**  $q^* = p^*$ .

**Remark 1.10.** • *Weak duality always holds. However, strong duality needs more well conditions.*

- *Let us discuss the following fact first:*

$$\sup_{\lambda \geq 0} \{f_0(\mathbf{x}) + \sum_i \lambda_i f_i(\mathbf{x})\} = \begin{cases} f_0(\mathbf{x}), & f_i(\mathbf{x}) \leq 0, i = 1, \dots, m \\ \infty, & \text{otherwise.} \end{cases}$$

Thus, we have

$$p^* = \inf_{\mathbf{x}} \sup_{\lambda \geq 0} L(\mathbf{x}, \lambda),$$

$$q^* = \sup_{\lambda \geq 0} \inf_{\mathbf{x}} L(\mathbf{x}, \lambda).$$

Therefore, the weak duality implies that

$$\sup_{\lambda \geq 0} \inf_{\mathbf{x}} L(\mathbf{x}, \lambda) \leq \inf_{\mathbf{x}} \sup_{\lambda \geq 0} L(\mathbf{x}, \lambda).$$

**Definition 1.11.** We refer to a pair  $(\bar{\mathbf{x}}, \bar{\mathbf{y}})$  as a *saddle-point* for  $f$  if

$$f(\bar{\mathbf{x}}, \mathbf{y}) \leq f(\bar{\mathbf{x}}, \bar{\mathbf{y}}) \leq f(\mathbf{x}, \bar{\mathbf{y}}),$$

for all  $(\mathbf{x}, \mathbf{y}) \in \text{dom}(f)$ . In other words,  $\bar{\mathbf{x}}$  minimizes  $f(\mathbf{x}, \bar{\mathbf{y}})$  and  $\bar{\mathbf{y}}$  minimizes  $f(\bar{\mathbf{x}}, \mathbf{y})$ . Saddle-point problems play an important role in **Game Theory and Generative Adversarial Networks**.

**Example 1.12.**

$$\begin{aligned} \min \quad & \|\mathbf{x}\|^2, \\ s.t. \quad & A\mathbf{x} = \mathbf{b}. \end{aligned}$$

- Lagrangian:  $L(\mathbf{x}, \nu) = \|\mathbf{x}\|^2 + \nu^\top (A\mathbf{x} - \mathbf{b})$ .

- Lagrange Dual Function:  $g(\boldsymbol{\nu}) = \inf_{\mathbf{x}} L(\mathbf{x}, \boldsymbol{\nu})$ . We know that  $\nabla_{\mathbf{x}} L(\mathbf{x}, \boldsymbol{\nu}) = 2\mathbf{x} + A^T \boldsymbol{\nu} = 0$ , thus  $\mathbf{x}^* = -\frac{1}{2} A^T \boldsymbol{\nu}$ . Take  $\mathbf{x}^*$  into Lagrangian, we obtain the Lagrange dual function

$$g(\boldsymbol{\nu}) = -\frac{1}{4} \boldsymbol{\nu}^T A A^T \boldsymbol{\nu} - \boldsymbol{\nu}^T \mathbf{b}$$

- Dual problem:  $\max -\frac{1}{4} \boldsymbol{\nu}^T A A^T \boldsymbol{\nu} - \boldsymbol{\nu}^T \mathbf{b}$ .
- Weak duality:

$$\sup_{\boldsymbol{\nu}} \left\{ -\frac{1}{4} \boldsymbol{\nu}^T A A^T \boldsymbol{\nu} - \boldsymbol{\nu}^T \mathbf{b} \right\} \leq \min_{\mathbf{x}} \{ \|\mathbf{x}\|^2 \mid A\mathbf{x} = \mathbf{b} \}.$$

**Example 1.13.** (Linear Programming) Recall the example of transportation problem in OM.

$$\begin{aligned} \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x}, \\ \text{s.t.} \quad & A\mathbf{x} = \mathbf{b}, \\ & \mathbf{x} \succeq 0. \end{aligned}$$

- Lagrangian:

$$L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = \mathbf{c}^T \mathbf{x} - \boldsymbol{\lambda}^T \mathbf{x} + \boldsymbol{\nu}^T (A\mathbf{x} - \mathbf{b}) = (\mathbf{c} - \boldsymbol{\lambda} + A^T \boldsymbol{\nu})^T \mathbf{x} - \boldsymbol{\nu}^T \mathbf{b}.$$

- Lagrange Dual Function:

$$g(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \begin{cases} -\boldsymbol{\nu}^T \mathbf{b}, & \mathbf{c} - \boldsymbol{\lambda} + A^T \boldsymbol{\nu} = 0, \\ -\infty, & \text{otherwise.} \end{cases}$$

- Dual problem:

$$\begin{aligned} \max_{\boldsymbol{\lambda}, \boldsymbol{\nu}} \quad & -\boldsymbol{\nu}^T \mathbf{b}, \\ \text{s.t.} \quad & \mathbf{c} - \boldsymbol{\lambda} + A^T \boldsymbol{\nu} = 0, \\ & \boldsymbol{\lambda} \succeq 0. \end{aligned}$$

This is equivalent to

$$\begin{aligned} \min_{\boldsymbol{\nu}} \quad & \boldsymbol{\nu}^T \mathbf{b}, \\ \text{s.t.} \quad & \mathbf{c} + A^T \boldsymbol{\nu} \succeq 0. \end{aligned}$$

**Example 1.14.**

$$\begin{aligned} \min_{\mathbf{x}} \quad & \|\mathbf{x}\|, \\ \text{s.t.} \quad & A\mathbf{x} = \mathbf{b}. \end{aligned}$$

It seems that we cannot obtain the Lagrange dual function via the directly derivation. How to do? We will learn and adapt conjugate function to handle this problem.